

Badger: Fast Key-Value DB in Go

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Dgraph Labs

- Fast, Distributed graph database.
- Sparse data sets.
- Lots of relationships.

https://dgraph.io



What is Badger?



- Badger is an embedded key-value database, written in Go.
- Licensed under Apache 2.0.

go get github.com/dgraph-io/badger/...

Current Status

- Closing v2.0.
- Close to 3500 Github starts.
- 42 contributors.
- Used by Dgraph, Go-IPFS, 0-stor, Sandglass.





Basic Operations



Set a key-value



```
func set() error {
    fmt.Println("\nRunning SET")
    return db.Update(func(txn *badger.Txn) error {
        if err := txn.Set([]byte("foo"), []byte("bar")); err != nil {
            return err
        }
        fmt.Println("Set foo to bar")
        return nil
    })
}
```

Get a key-value



```
func get() error {
    fmt.Println("\nRunning GET")
    return db.View(func(txn *badger.Txn) error {
        item, err := txn.Get([]byte("foo")) // handle err
        if err != nil {
            return err
        }
        val, err := item.Value() // handle err
        if err != nil {
            return err
        }
        fmt.Printf("The value is: %s\n", val)
        return nil
    })
}
```

Iterate key-values



```
func iterate() error {
    fmt.Println("\nRunning ITERATE")
    return db.View(func(txn *badger.Txn) error {
        opts := badger.DefaultIteratorOptions
        it := txn.NewIterator(opts)
        defer it.Close()
        for it.Rewind(); it.Valid(); it.Next() {
            k := it.Item().Key()
            v, err := it.Item().Value() // handle err
            if err != nil {
                return err
            }
            fmt.Printf("key=%s, value=%s\n", k, v)
        }
        return nil
   })
}
```

Run the code



Run

```
func main() {
    opt := badger.DefaultOptions
    opt.Dir = "/tmp/db"
    opt.ValueDir = opt.Dir
   var err error
   db, err = badger.Open(opt)
    if err != nil {
        panic(err)
    }
    defer db.Close()
    fmt.Println("DB opened")
    set()
   get()
    iterate()
    fmt.Println("DB done")
}
```

Badger != replacement for Go map



Motivation and Outcome



Cgo is not Go

Some people, when confronted with a problem, think "I know, I'll use cgo."

Now they have two problems.

->Cgo is not Go, Dave Cheney



RocksDB

- Great write throughput.
- Okay read throughput.

Cons:

• Required Cgo.



BoltDB

- Pure Go.
- Great read throughput.

Cons:

• Bad write throughput.

Why build it?

- Go native key-value DB for Dgraph.
- No compromise in read-write performance.
- Avoid Cgo.



What did we spend?

- Spent a few months.
- Built with <1 full-time gopher.
- Aka, the power of Go!





What did we gain?

- A faster key-value DB for Go.
- Ability to run Go profilers all the way down to disk.
- Clean Go code (no C).



Launch Reception

- Within 12 hours of blog post release
- First page of HN for a day.
- 355 points, 96 comments.
- 1250 Github stars in 4 days.









Recruiters loved it!



• Got 3 different emails from 3 different recruiters...

Recruiters loved it!

• For jobs in the same company.









Two common Trees

- LSM trees
- B+ trees



LSM Trees





Compaction continues creating fewer, larger and larger files

- More levels
- High write throughput
- High read latency
- Example: RocksDB

B+Trees





- Fewer levels
- Low write throughput
- Low read latency
- Example: BoltDB

Badger is based on LSM trees.



LSM Trees





Writes in LSM trees: Memtable to L0





Writes in LSM trees: L0 to L1





Writes in LSM trees: Li to Li+1





What makes Badger unique?

- Based on WiscKey paper by Uni Wisconsin-Madison.
- Separates keys from values.
- Stores keys in LSM tree.
- Stores value in value log.



Write to Value Log





• Write value, get pointer.

More keys per table





Smaller LSM tree





Typical Badger setup




Advantages of smaller LSM tree

- Can be kept in RAM.
- Low read amplification (fewer lookups).
- Low write amplification (fewer compactions).
- \propto Number of keys.



Usenet Express

- Hundreds of terabytes of data.
- Few gigabytes of LSM tree.



Reads in Badger: LSM tree





Reads in Badger: Bloom Filters





Reads in Badger: Value Log





• Once key found in LSM tree, read from value log.

Badger is FAST-er



Data Loading: Badger vs Go-RocksDB





Data Loading Performance

• As value size increases, Badger's becomes 11.7x faster.

Data Loading: Badger vs BoltDB





Data Loading Performance

• 11x - 22x faster than BoltDB on all value sizes.

Random Reads: Badger vs Go-RocksDB





Random Read Latency

Random Get latency is 3.7x - 5.3x lower than RocksDB. •

Random Reads: Badger vs BoltDB





• Random Get latency is slightly better or worse, depending on value size.

Various other benchmarks

Range iteration latency, etc.

Can be found on https://blog.dgraph.io/

Benchmarking code is open sourced.

github.com/dgraph-io/badger-bench (https://github.com/dgraph-io/badger-bench)



Features



Concurrent Transactions



Badger uses Oracle to achieve concurrent lock-free transactions.



Concurrent Writes

Batch up writes from multiple transactions. Amortize cost of disk write.

No wait -> Smart Batching.



Smart Batching in Go



```
reqs := make([]*request, 0, 10)
                                          closedCase:
for {
                                              close(db.writeCh)
                                              for r := range db.writeCh {
    var r *request
    select {
                                                   reqs = append(reqs, r)
                                              }
    case r = <-db.writeCh:</pre>
    case <-lc.HasBeenClosed():</pre>
         goto closedCase
                                              pendingCh <- struct{}{}</pre>
    }
                                              writeRequests(reqs)
                                              return
    for {
         reqs = append(reqs, r)
                                          writeCase:
         reqLen.Set(int64(len(reqs)))
                                              go writeRequests(reqs)
                                              reqs = make([]*request, 0, 10)
         select {
                                              reqLen.Set(0)
         case r = <-db.writeCh:</pre>
         case pendingCh <- struct{}{}:</pre>
             goto writeCase
         case <-lc.HasBeenClosed():</pre>
             goto closedCase
         }
```

Multi Version Concurrency Control





Badger stores multiple versions of the key.

Provides direct access to the versions, via iterate.

Crash Resilience

LSM Memtables can be lost to crashes.

Can be recovered from value log on restart.



Value Log Garbage Collection



Why?

Value log would keep growing with every Set.

Older versions of keys can be deleted.

Corresponding values can be deleted from value log.



Stage 1: Punch Holes (v2.0 in Linux)





Punch Holes



Stage 2: Move entries, Delete log





Dealing with Qu-err-key file systems.





Would a file delete reclaim space in the f is system?

Delete, no reclaim



No

```
if err := t.fd.Truncate(0); err != nil {
    // This is very important to let the FS know
    // that the file is deleted.
    return err
}
```

Truncate the file before deleting.

Would closing a file sync its contents to d



Close, no-sync



No

```
if err := lf.fd.Sync(); err != nil {
    return errors.Wrapf(err, "Unable to sync value log: %q", lf.path)
}
if err := lf.fd.Close(); err != nil {
    return errors.Wrapf(err, "Unable to close value log: %q", lf.path)
}
```

Explicitly sync file before closing.

Can a new synced file be lost?



Create, no-found



Yes

```
f, err := os.Open(dir)
if err != nil {
    return errors.Wrapf(err, "While opening directory: %s.", dir)
}
err = f.Sync()
closeErr := f.Close()
if err != nil {
    return errors.Wrapf(err, "While syncing directory: %s.", dir)
}
return errors.Wrapf(closeErr, "While closing directory: %s.", dir)
```

Sync a directory just like you would sync a file.



Crash, no-clean

Yes.

Add checksums to know when to truncate a file.



Who should use Badger?



Don't use Badger if...

You no Go! (C++, Java)

You have a single-threaded sequential workload.

You have a small, read-only workload.

All your data can fit in memory easily.



Use Badger if...



You Go!

- You want to avoid Cgo.
- You want a performant read-write workload.
- You access data concurrently (many goroutines accessing data).
- You need 3-dimensional access.

Future Work

Encryption at rest.

Others? (tell us what you need)

Work on Badger and Dgraph. Come join us!

github.com/dgraph-io/badger (https://github.com/dgraph-io/badger)

github.com/dgraph-io/dgraph(https://github.com/dgraph-io/dgraph)

Careers at Dgraph (https://dgraph.io/about.html)



Talk to us on Wechat


Thank you

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